

A game-theoretical model for green urban logistics coordination: exact and data-driven methods

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Motivation

URBANE: Upscaling Innovative Green Urban Logistics Solutions Through Multi-Actor Collaboration and PI-Inspired Last Mile Deliveries.
<https://www.urbane-horizoneurope.eu/>

- A 3.5-year Horizon Europe project (2022-2026) on novel, **sustainable**, safe, resilient and **effective last-mile delivery solutions**, combining green automated vehicles and shared space utilisation models.
- **Living Labs:**
Helsinki (FI), Bologna (IT), Valladolid (ES), Thessaloniki (GR)
- **Twinning Living Labs:**
Barcelona (ES), Karlsruhe (DE)
- **Follower Cities:**
Aarhus (DK), Antwerp (BE), La Rochelle (FR), Mechelen (BE), Prague (CZ), Ravenna (IT)

Motivation - Challenges of Bologna LL (1)

- Coordinate the activities of the logistics operators
- Reduction of the distance travelled by the delivery vehicles and consequently reduction of the negative externalities produced by them (road congestion, noise, CO_2 and pollutant emissions, etc.)

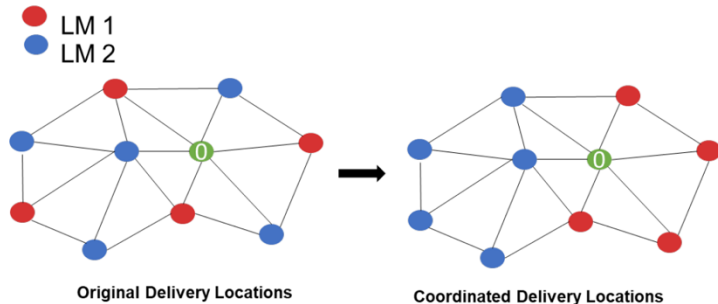


Figure: Solution Example (Last Mile only)

Motivation - Challenges of Bologna LL (2)



Figure: Bologna LL Limited Traffic/Emission Zone (LTZ/LEZ)

Bilevel Optimization for Greener Last Mile Logistics (1)

The goal is to consolidate parcel deliveries outside the LTZ/LEZ, and reduce the overlap of the operating areas of the Delivery Service Providers (DSPs) in the LEZ

Two-echelon (2E) multi-commodity routing model:

- **First echelon:** First Mile (FM) DSPs which deliver parcels to satellite hubs at the edges of the LEZ
- **Second echelon:** Last Mile (LM) DSPs must pickup parcels at satellite hubs and deliver them to customers

Several non-cooperative agents need to be coordinated $\implies$ Bilevel Optimization (BO) [Dempe, 2002]

Bilevel Optimization for Greener Last Mile Logistics (2)

BO approach:

Stackelberg game where the leader represents a regulator whose goal is to minimize emissions in the LEZ

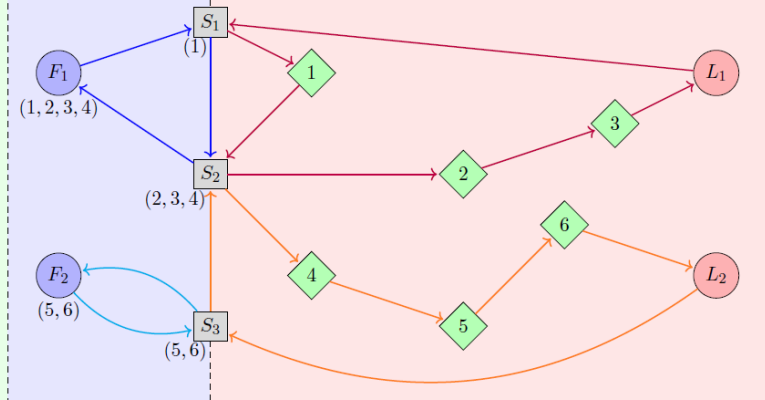
- The **leader** assigns parcels to satellite hubs and to LM DSPs. The followers seek to minimize their own routing costs and maximize client satisfaction
- Each **FM follower** solve a Vehicle Routing Problem with Time Windows (VRPTW)
- Each **LM follower** solves a pickup-and-delivery VRPTW (PDPTW)

⇒ **Single leader, multi-follower bilevel optimization formulation** for 2E multi-commodity routing problem to facilitate the city's goals of reducing emissions in its LEZ

Leader: assigns parcels to satellites and to LMs

FM followers (VRPTW)

LM followers (PDPTW)



FM depot node with its parcels



Satellite node with its assigned parcels



LM depot node



Parcel destination (customer) node

Leader Problem

Minimize emissions in both echelons:

$$\min_{\lambda, \mathbf{y}} \quad \sum_{k \in \mathcal{K}_1} \sum_{(i,j) \in \mathcal{A}_1} E_{ij}^k x_{ij,f}^k + \sum_{k \in \mathcal{K}_2} \sum_{(i,j) \in \mathcal{A}_2} E_{ij}^k x_{ij,d}^k$$

$$\text{s.t.} \quad \sum_{f \in \mathcal{F}} \sum_{p \in \mathcal{P}_f} \lambda_p^s \leq C_s \quad \forall s \in \mathcal{S}$$

$$\sum_{s \in \mathcal{S}} \lambda_p^s = 1 \quad \forall p \in \mathcal{P}$$

$$\sum_{d \in \mathcal{D}} y_p^d = 1 \quad \forall p \in \mathcal{P}$$

$$G(\lambda, \mathbf{y}) \geq 0$$

$$\lambda_p^s \in \{0, 1\} \quad \forall p \in \mathcal{P}, s \in \mathcal{S}$$

$$y_p^d \in \{0, 1\} \quad \forall p \in \mathcal{P}, d \in \mathcal{D}$$

$$(\mathbf{x}_f, \mathbf{t}_f) \in \text{VRPTW}(\lambda_f) \quad \forall f \in \mathcal{F}$$

$$(\mathbf{x}_d, \mathbf{t}_d) \in \text{PDPTW}(\lambda, \mathbf{y}_d) \quad \forall d \in \mathcal{D}$$

First Mile Follower Problem

Each FM follower $f \in \mathcal{F}$ minimizes the total route cost associated with delivering its set of parcels $\mathcal{P}_f$

$$\begin{aligned}
 \text{VRPTW}(\lambda_f) : \min_{\mathbf{x}_f, \mathbf{t}_f} \quad & \sum_{k \in \mathcal{K}_f} \sum_{(i,j) \in \mathcal{A}_1} c_{ij} x_{ij,f}^k \\
 \text{s.t.} \quad & \sum_{i \in \mathcal{V}_1} \sum_{k \in \mathcal{K}_f} x_{is,f}^k \geq \lambda_p^s & \forall p \in \mathcal{P}_f, s \in \mathcal{S} \\
 & \sum_{s \in \mathcal{S}} x_{D_f s, f}^k \leq 1 & \forall k \in \mathcal{K}_f \\
 & \sum_{i \in \mathcal{V}_1} x_{ij, f}^k - \sum_{i \in \mathcal{V}_1} x_{ji, f}^k = 0 & \forall k \in \mathcal{K}_f, j \in \mathcal{V}_1 \\
 & t_{i, f}^k + T_{ij} \leq t_{j, f}^k + (1 - x_{ij, f}^k) M_{ij} & \forall k \in \mathcal{K}_f, (i, j) \in \mathcal{A}_1 \\
 & e_i \sum_{j \in \mathcal{V}_1} x_{ij, f}^k \leq t_{i, f}^k \leq l_i \sum_{j \in \mathcal{V}_1} x_{ij, f}^k & \forall k \in \mathcal{K}_f, i \in \mathcal{V}_1 \\
 & x_{ij, f}^k \in \{0, 1\} & \forall k \in \mathcal{K}_f, (i, j) \in \mathcal{A}_1 \\
 & t_{i, f}^k \geq 0 & \forall k \in \mathcal{K}_f, i \in \mathcal{V}_1
 \end{aligned}$$

Last Mile Follower Problem (1)

Each LM follower $d \in \mathcal{D}$ minimizes its total routing cost, as well as time window violations

$$\begin{aligned} \text{PDPTW}(\lambda, \mathbf{y}_d) : \min_{\mathbf{x}_d, \mathbf{t}_d} & \sum_{k \in \mathcal{K}_d} \sum_{(i,j) \in \mathcal{A}_2} c_{ij} x_{ij,d}^k \\ & + \pi^L \sum_{k \in \mathcal{K}_d} \sum_{i \in \mathcal{V}_2} \max\{t_{i,d}^k - l_i \sum_{j \in \mathcal{V}_2} x_{ij,d}^k, 0\} \\ & + \pi^E \sum_{k \in \mathcal{K}_d} \sum_{i \in \mathcal{V}_2} \max\{e_i \sum_{j \in \mathcal{V}_2} x_{ij,d}^k - t_{i,d}^k, 0\} \end{aligned}$$

Last Mile Follower Problem (2)

$$\begin{aligned}
 \text{s.t.} \quad & \sum_{j \in \mathcal{V}_2} x_{j\delta(p),d}^k = y_p^d && \forall p \in \mathcal{P}, \forall k \in \mathcal{K}_d \\
 & \sum_{s \in \mathcal{S}} \sum_{j \in \mathcal{V}_2} \lambda_p^s x_{sj,d}^k \geq y_p^d && \forall p \in \mathcal{P}, \forall k \in \mathcal{K}_d \\
 & \sum_{s \in \mathcal{S}} x_{D_d s,d}^k \leq 1 && \forall k \in \mathcal{K}_d \\
 & \sum_{j \in \mathcal{V}_2} x_{sj,d}^k \leq 1 && \forall k \in \mathcal{K}_d, s \in \mathcal{S} \\
 & \sum_{i \in \mathcal{V}_2} x_{ij,d}^k - \sum_{i \in \mathcal{V}_2} x_{ji,d}^k = 0 && \forall k \in \mathcal{K}_d, j \in \mathcal{V}_2 \\
 & \sum_{s \in \mathcal{S}} \lambda_p^s (t_{s,d}^k + T_{s\delta(p)}) \leq t_{\delta(p),d}^k + (1 - y_p^d) N_p^k && \forall p \in \mathcal{P}, k \in \mathcal{K}_d \\
 & t_{i,d}^k + x_{ij,d}^k T_{ij} \leq t_{j,d}^k + (1 - x_{ij,d}^k) M_{ij} && \forall k \in \mathcal{K}_d, (i,j) \in \mathcal{A}_2 \\
 & x_{ij,d}^k \in \{0, 1\} && \forall (i,j) \in \mathcal{A}_2, k \in \mathcal{K}_d \\
 & t_{i,d}^k \geq 0 && \forall i \in \mathcal{V}_2, k \in \mathcal{K}_d
 \end{aligned}$$

Solution Approach 1 — Cutting Plane Algorithm

Initialization: relax followers' best-response constraints in Leader problem
 $\implies$ High-Point Relaxation (HPR) and reformulate HPR as MILP (max and bilinear term linearizations)

At iteration n :

- ❶ Solve HPR to obtain lower bound (LB), candidate (λ^n, y^n) solution and estimates of follower's objective function values (OFV): $\hat{V}_f^n, \hat{V}_d^n$
- ❷ Solve (λ^n, y^n) -parameterized follower problems to obtain upper bound (UB) and best-responses OFVs: V_f^n, V_d^n
- ❸ **Optimality cuts generation:**
 - if $\hat{V}_f^n - V_f^n > \epsilon$ generate FM optimality cut for FM f
 - if $\hat{V}_d^n - V_d^n > \epsilon$ generate LM optimality cut for LM d
- ❹ Add **interdiction cut**

Exact cutting plane algorithm that terminates if LB and UB converge, or if no optimality cuts are added

Solution Approach 2 — Constraint Learning (CL)

In BO, we work for the leader: follower problems are the leader's representations of FM and LM DSPs' operations

Motivation for a CL heuristic approach:

- Leader decisions are **parcel hub and LM assignments**, precise DSP routes are not the target output
- Data-driven **prediction** of FM and LM best responses may be sufficient to find efficient leader decisions
- **Training** predictive models for learning FM and LM best responses consists of solving multiple parameterized VRPTWs and PDPTWs for which advanced algorithms are available in the literature

CL-Based Leader Problem

$$\begin{aligned}
 & \min_{\boldsymbol{\lambda}, \mathbf{y}, \boldsymbol{\gamma}} && \sum_{f \in \mathcal{F}} E_f + \sum_{d \in \mathcal{D}} E_d + \rho_S \sum_{s \in \mathcal{S}} \gamma_s + \rho_D \sum_{d \in \mathcal{D}} \gamma_d \\
 & \text{s.t.} && \sum_{f \in \mathcal{F}} \sum_{p \in \mathcal{P}_f} \lambda_p^s \leq C_s && \forall s \in \mathcal{S} \\
 & && \sum_{s \in \mathcal{S}} \lambda_p^s = 1 && \forall p \in \mathcal{P} \\
 & && \sum_{d \in \mathcal{D}} y_p^d = 1 && \forall p \in \mathcal{P} \\
 & && E_f = \text{Pred}(\boldsymbol{\lambda}_f) && \forall f \in \mathcal{F} \\
 & && E_d = \text{Pred}(\boldsymbol{\lambda}, \mathbf{y}_d) && \forall d \in \mathcal{D} \\
 & && \gamma_s \geq \lambda_p^s && \forall p \in \mathcal{P}, s \in \mathcal{S} \\
 & && \gamma_d \geq y_p^d && \forall p \in \mathcal{P}, d \in \mathcal{D} \\
 & && \lambda_p^s \in \{0, 1\} && \forall p \in \mathcal{P}, s \in \mathcal{S} \\
 & && y_p^d \in \{0, 1\} && \forall p \in \mathcal{P}, d \in \mathcal{D}
 \end{aligned}$$

$\text{Pred}(\cdot)$ constraints can be formulated using MILP [[Maragno et al., 2023](#)]

ρ_S and ρ_D are design penalties

Numerical Experiments — Setup

All models are implemented in Python:

- Exact approach implemented using CPLEX 22.1
- Data-driven approach implemented using Gurobi and `gurobi-machinelearning 1.4.0`

We compared 3 approaches:

- Uncoordinated deliveries (heuristic (λ, y) assignment)
- Fully Coordinated deliveries (Bilevel: Exact)
- Fully Coordinated deliveries (Bilevel: Data-Driven)

Numerical Results — Unassigned Deliveries

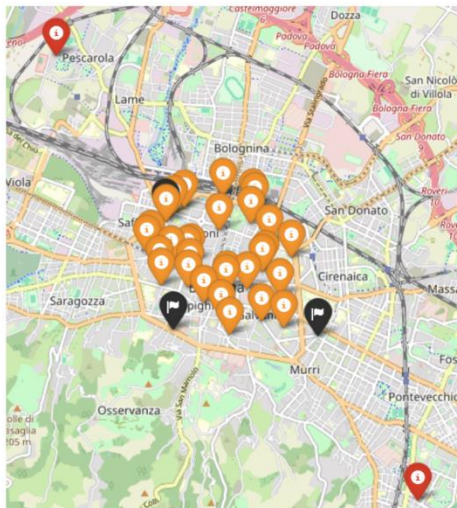


Figure: Problem instance with $|\mathcal{D}| = 2$, $|\mathcal{S}| = 3$ & $|\mathcal{P}| = 30$

Numerical Results — Assigned and Coordinated Deliveries



Figure: Coordinated approach leads to consolidation of delivery locations (berlin_2_2_3_30) (only last mile solution shown)

Numerical Results — Optimization Methods

Table: Performance of different modelling approaches on the Madrid instances of the bilevel 2E problem. Iter. denotes the number of iterations; Time is the runtime in seconds; Obj. is the objective value of the best feasible solution found; Gap is the optimality gap upon termination in percentage.

Method	Exact				Data-Driven		Uncoordinated	
Instance	Obj.	Time	Iter.	Gap	Obj.	Time	Obj.	Time
madrid_2.2.2.10	405.97	50.23	3	0.00	405.97	0.48	656.76	0.40
madrid_2.2.2.30	828.88	4225.72	4	37.06	824.11	306.49	1177.42	303.13
madrid_2.2.2.50	1309.35	3916.70	4	51.17	1283.01	348.77	1684.25	304.78
madrid_2.2.3.10	431.24	705.12	5	0.00	438.29	0.64	628.00	0.59
madrid_2.2.3.30	794.48	4279.60	4	40.36	842.87	344.56	1169.94	302.16
madrid_2.2.3.50	1414.53	3936.97	4	58.09	1352.54	507.74	1656.34	309.48
madrid_2.3.2.10	405.97	144.48	4	0.00	460.10	1.34	760.02	1.25
madrid_2.3.2.30	852.85	3617.98	4	40.74	853.64	312.00	1360.28	604.82
madrid_2.3.2.50	1398.35	4283.50	4	56.05	1334.38	391.53	1514.09	613.17
madrid_2.3.3.10	431.24	3796.13	9	0.00	438.29	1.30	968.20	1.56
madrid_2.3.3.30	856.77	2709.28	4	46.41	844.46	472.39	1312.28	605.92
madrid_2.3.3.50	1329.87	4528.09	4	56.04	1224.56	550.30	2030.28	643.99

Numerical Results — Total Emissions

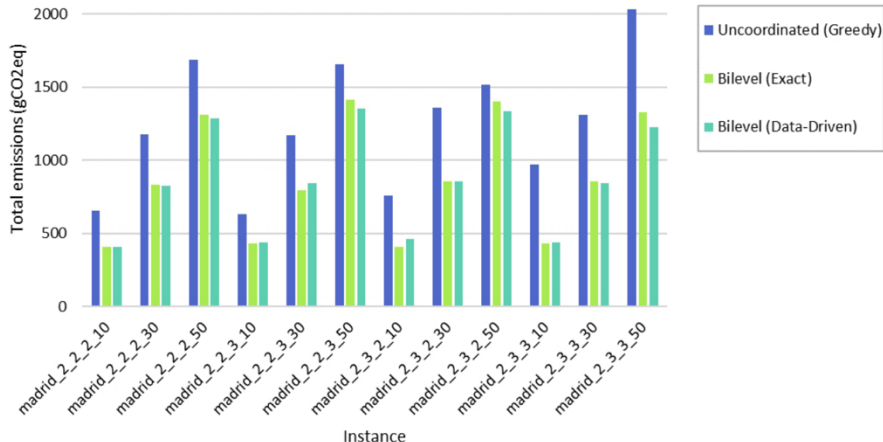


Figure: Total emissions (gCO₂eq) across all Madrid instances.

Summary and Conclusions

- New collaborative delivery model for greener urban logistics.
- Bilevel optimization approach where followers represent delivery service providers.
- Exact and data-driven solution methods are developed and implemented.
- Coordinated model is shown to strike a balance between carbon emissions reduction and delivery service providers operating costs.
- Constraint learning-based approach provides a tractable solution method for optimizing city-scale operations.
- Reduced daily operational emissions by 15%

References

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